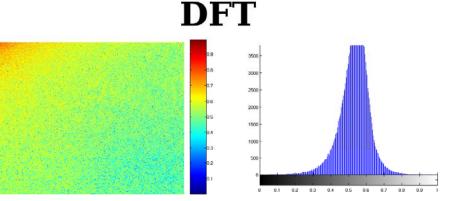
Nuno Vasconcelos UCSD

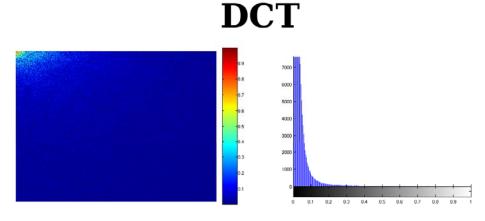
Discrete Fourier Transform

- last classes, we have studied the DFT
- due to its computational efficiency the DFT is very popular
- however, it has strong disadvantages for some applications
 - it is complex
 - it has poor energy compaction
- energy compaction
 - is the ability to pack the energy of the spatial sequence into as few frequency coefficients as possible
 - this is very important for image compression
 - we represent the signal in the frequency domain
 - if compaction is high, we only have to transmit a few coefficients
 - instead of the whole set of pixels

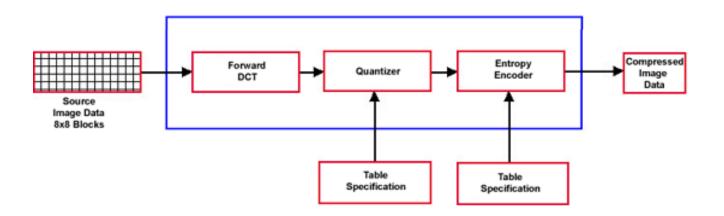
- a much better transform, from this point of view, is the DCT
 - in this example we see the amplitude spectra of the image above
 - under the DFT and DCT
 - note the much more concentrated histogram obtained with the DCT
- why is energy compaction important?
 - the main reason is image compression
 - turns out to be beneficial in other applications







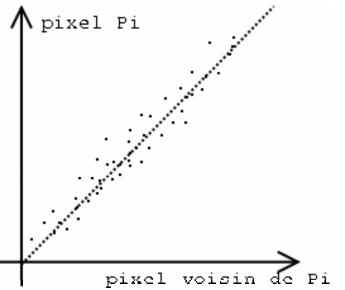
• an image compression system has three main blocks



- a transform (usually DCT on 8x8 blocks)
- a quantizer
- a lossless (entropy) coder
- each tries to throw away information which is not essential to understand the image, but costs bits

- the transform throws away correlations
 - if you make a plot of the value of a pixel as a function of one of its neighbors





- you will see that the pixels are highly correlated (i.e. most of the time they are very similar)
- this is just a consequence of the fact that surfaces are smooth

- the transform eliminates these correlations
 - this is best seen by considering the 2-pt transform
 - note that the first coefficient is always the DC-value

X[0] = x[0] + x[1]

- an orthogonal transform can be written in matrix form as

$$X = Tx, \qquad T^T T = I$$

- i.e. T has orthogonal columns
- this means that

$$X[1] = x[0] - x[1]$$

- note that if x[0] similar to x[1], then

$$\begin{cases} X[0] = x[0] + x[1] \approx 2x[0] \\ X[1] = x[0] - x[1] \approx 0 \end{cases}$$

- the transform eliminates these correlations
 - note that if x[0] similar to x[1], the

$$\begin{cases} X[0] = x[0] + x[1] \approx 2x[0] \\ X[1] = x[0] - x[1] \approx 0 \end{cases}$$

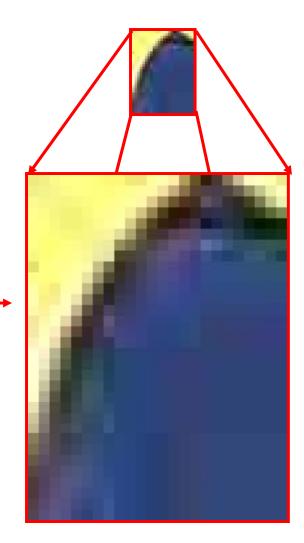
- in the transform domain we only have to transmit one number without any significant cost in image quality
- by "decorrelating" the signal we reduced the bit rate to $\frac{1}{2}$!
- note that an orthogonal matrix

$$T^T T = I$$

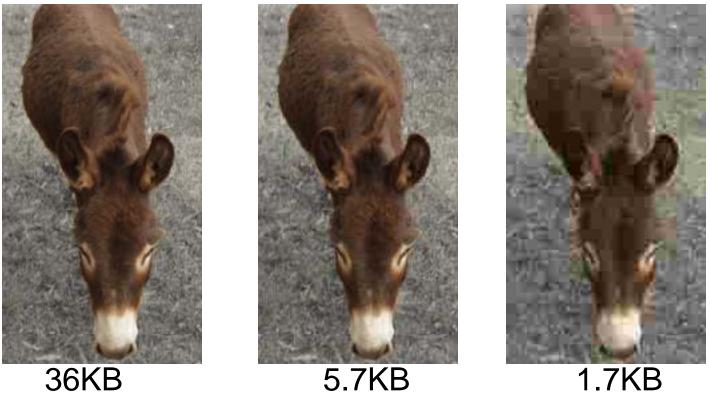
applies a rotation to the pixel space

- this aligns the data with the canonical axes

- a second advantage of working in the frequency domain
 - is that our visual system is less sensitive to distortion around edges
 - the transition associated with the edge masks our ability to perceive the noise
 - e.g. if you blow up a compressed picture, it is likely to look like this
 - in general, the compression errors are more annoying in the smooth image regions



• three JPEG examples



- note that the blockiness is more visible in the torso

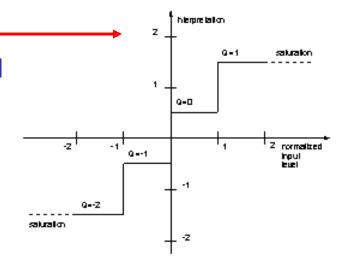
- important point: by itself, the transform
 - does not save any bits
 - does not introduce any distortion
- both of these happen when we throw away information
- this is called "lossy compression" and implemented by the quantizer
- what is a quantizer?
 - think of the round() function, that rounds to the nearest integer
 - round(1) = 1; round(0.55543) = 1; round (0.0000005) = 0
 - instead of an infinite range between 0 and 1 (infinite number of bits to transmit)
 - the output is zero or one (1 bit)
 - we threw away all the stuff in between, but saved a lot of bits
 - a quantizer does this less drastically

- it is a function of this type
 - inputs in a given range are mapped to the same output
- to implement this, we
 - 1) define a quantizer step size Q
 - 2) apply a rounding function

$$x_q = round\left(\frac{x}{Q}\right)$$



- the larger the Q, the less reconstruction levels we have
- more compression at the cost of larger distortion
- e.g. for x in [0,255], we need 8 bits and have 256 color values
- with Q = 64, we only have 4 levels and only need 2 bits



- note that we can quantize some frequency coefficients more heavily than others by simply increasing Q
- this leads to the idea of a quantization matrix
- we start with an image block (e.g. 8x8 pixels)

						52	55	61	66	70	61	64	73
						63	59	55	90	109	85	69	72
						62	59	68	113	144	104	66	73
						63	58	71	122	154	106	70	69
					\longleftrightarrow	67	61	68	104	126	88	68	70
						79	65	60	70	77	68	58	75
										55			
			_	-		87	79	69	68	65	76	78	94
						-							-

0

 $^{-1}$

0

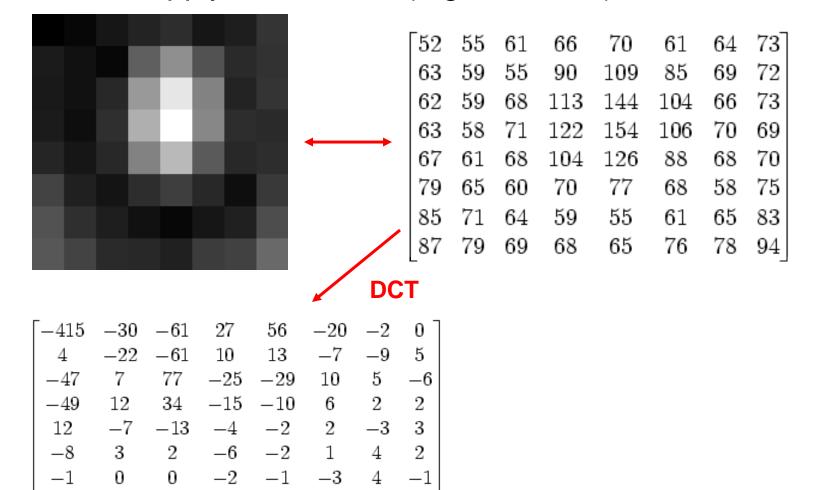
-4

 $^{-1}$

0

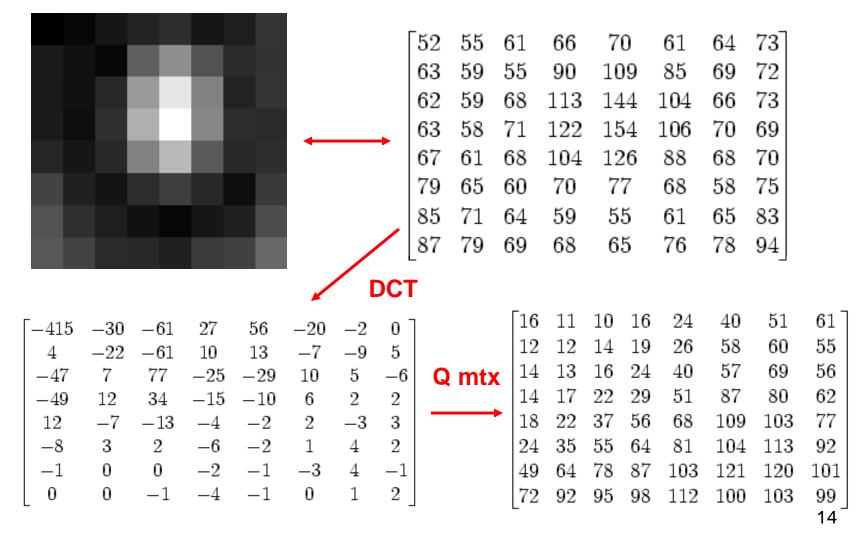
1

next we apply a transform (e.g. 8x8 DCT)



2

and quantize with a varying Q



note that higher frequencies are quantized more heavily

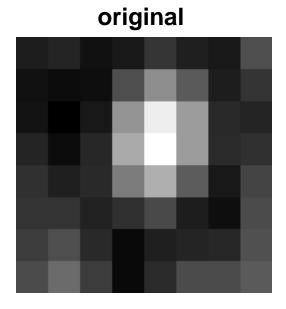
5161] 16 11 10 16 24 40Q mtx 12 12 14 1 increasing frequency $13 \ 16 \ 24 \ 40$ 5714 69 5617 22 29 51 87 14 80 6218 22 37 56 68 109 103 77 24 35 55 64 81 104 113 9249 64 78 87 103 121 120 10172 92 95 98 112 100 10399

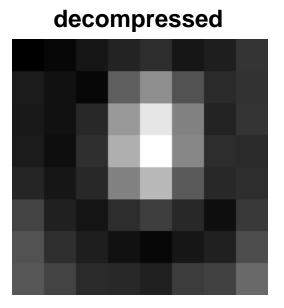
 in result, many high frequency coefficients are simply wiped out **Quantized DCT**

 this saves a lot of bits, but we no longer have an exact replica of original image block

	DC	T								q	uan	tize	d D	СТ				
[-41	5 -	-30	-61	27	56	-20	-2	0		-26	-3	$^{-6}$	2	2	$^{-1}$	0	0	
4	-	-22	-61	10	13	-7	$^{-9}$	5		0	-2	-4	1	1	0	0	0	
-4'	7	7	77	-25	-29	10	5	$^{-6}$		-3	1	5	$^{-1}$	$^{-1}$	0	0	0	
-49	9	12	34	-15	-10	6	2	2		-4	1	2	$^{-1}$	0	0	0	0	
12		-7	-13	-4	-2	2	$^{-3}$	3		1	0	0	0	0	0	0	0	
-8		3	2	$^{-6}$	-2	1	4	2		0	0	0	0	0	0	0	0	
-1		0	0	-2	-1	-3	4	-1		0	0	0	0	0	0	0	0	
0		0	$^{-1}$	-4	$^{-1}$	0	1	2		0	0	0	0	0	0	0	0	
inverse DCT original pixels																		
60	63	55	58	70	61	58	80			52	55 (61	66	70	61	64	1 7	3
58	56	56	83	108	88	63	71			63	59 3	55	90	109	85	69	9 7	2
60	52	62	113	150	116	70	67			62	59 (68 1	113	144	104	66	3 73	3
66	56	68	122	156	116	69	72			63	58 '	71 1	122	154	106	70	6	9
69	62	65	100	120	86	59	76		\neq	67	61 (68 1	104	126	88	68	3 7	0
68	68	61	68	78	60	53	78		7	79	65 (70	77	68	58		
74	82	67	54	63	64	65	83			85			59	55	61	65		
83	96	77	56	70	83	83	89			87			68	65	76	78		
Loo	00	•••	50		00	00	<u> </u>				19 (59	08	05	70	10	5 9	*_

 note, however, that visually the blocks are not very different





- we have saved lots of bits without much "perceptual" loss
- this is the reason why JPEG and MPEG work

• three JPEG examples









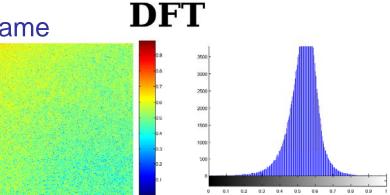
36KB

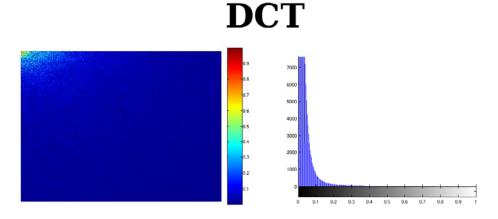


- note that the two images on the left look identical
- JPEG requires 6x less bits

- note that
 - the better the energy compaction
 - the larger the number of coefficients that get wiped out
 - the greater the bit savings for the same loss
- this is why the DCT is important
- we will do mostly the 1D-DCT
 - the formulas are simpler the insights the same
 - as always, extension to 2D is trivial







- the first thing to note is that there are various versions of the DCT
 - these are usually known as DCT-I to DCT-IV
 - they vary in minor details
 - the most popular is the DCT-II, also known as even symmetric DCT, or as "the DCT"

$$w[k] = \begin{cases} \frac{1}{2}, & k = 0\\ 1, & 1 \le k < N \end{cases}$$

$$C_{x}[k] = \begin{cases} \sum_{n=0}^{N-1} 2x[n] \cos\left(\frac{\pi}{2N}k(2n+1)\right), & 0 \le k < N\\ 0 & otherwise \end{cases}$$
$$x[n] = \begin{cases} \frac{1}{N} \sum_{k=0}^{N-1} w[k] C_{x}[k] \cos\left(\frac{\pi}{2N}k(2n+1)\right) & 0 \le n < N\\ 0 & otherwise \end{cases}$$

$$C_{x}[k] = \begin{cases} \sum_{n=0}^{N-1} 2x[n] \cos\left(\frac{\pi}{2N}k(2n+1)\right), & 0 \le k < N\\ 0 & otherwise \end{cases}$$

- from this equation we can immediately see that the DCT coefficients are real
- to understand the better energy compaction
 - it is interesting to compare the DCT to the DFT
 - it turns out that there is a simple relationship
- we consider a sequence x[n] which is zero outside of {0, ..., N-1}
- to relate DCT to DFT we need three steps

• step 1): create a sequence

$$y[n] = x[n] + x[2N - n - 1]$$

=
$$\begin{cases} x[n], & 0 \le n < N \\ x[2N - n - 1], & N \le n < 2N \end{cases}$$

step 2): compute the 2N-point DFT of y[n]

$$Y[k] = \sum_{n=0}^{2N-1} y[n] e^{-j\frac{2\pi}{2N}kn}, \quad 0 \le k < 2N$$

• step 3): rewrite as a function of N terms only

$$Y[k] = \sum_{n=0}^{N-1} y[n]e^{-j\frac{2\pi}{2N}kn} + \sum_{n=N}^{2N-1} y[n]e^{-j\frac{2\pi}{2N}kn}$$

• step 3): rewrite as a function of N terms only

$$Y[k] = \sum_{n=0}^{N-1} x[n] e^{-j\frac{2\pi}{2N}kn} + \sum_{n=N}^{2N-1} x[2N-n-1] e^{-j\frac{2\pi}{2N}kn}$$
$$= (m = 2N - 1 - n, \quad n = 2N - 1 - m) =$$
$$= \sum_{n=0}^{N-1} x[n] e^{-j\frac{2\pi}{2N}kn} + \sum_{m=0}^{N-1} x[m] e^{-j\frac{2\pi}{2N}k(2N-1-m)}$$
$$= \sum_{n=0}^{N-1} x[n] \left\{ e^{-j\frac{2\pi}{2N}kn} + e^{j\frac{2\pi}{2N}kn} \underbrace{e^{-j\frac{2\pi}{2N}k(2N-1-m)}}_{1} e^{j\frac{2\pi}{2N}k} \right\}$$
$$= \sum_{n=0}^{N-1} x[n] \left\{ e^{-j\frac{2\pi}{2N}kn} + e^{j\frac{2\pi}{2N}kn} e^{j\frac{2\pi}{2N}k} \right\}$$

 to write as a cosine we need to make it into two "mirror" exponents

• step 3): rewrite as a function of N terms only

$$Y[k] = \sum_{n=0}^{N-1} x[n] \left\{ e^{-j\frac{2\pi}{2N}kn} + e^{j\frac{2\pi}{2N}kn} e^{j\frac{2\pi}{2N}k} \right\}$$
$$= \sum_{n=0}^{N-1} x[n] e^{j\frac{\pi}{2N}k} \left\{ e^{-j\frac{2\pi}{2N}kn} e^{-j\frac{\pi}{2N}k} + e^{j\frac{2\pi}{2N}kn} e^{j\frac{\pi}{2N}k} \right\}$$
$$= \sum_{n=0}^{N-1} 2x[n] e^{j\frac{\pi}{2N}k} \cos\left(\frac{\pi}{2N}k(2n+1)\right)$$
$$= e^{j\frac{\pi}{2N}k} \sum_{n=0}^{N-1} 2x[n] \cos\left(\frac{\pi}{2N}k(2n+1)\right)$$

– from which

$$Y[k] = e^{j\frac{\pi}{2N}k} C_x[k], \qquad 0 \le k < 2N$$

• it follows that

$$C_{x}[k] = \begin{cases} e^{-j\frac{\pi}{2N}k} Y[k], & 0 \le k < N \\ 0, & otherwise \end{cases}$$

• in summary, we have three steps

$$\underbrace{x[n]}_{N-pt} \leftrightarrow \underbrace{y[n]}_{2N-pt} \quad \stackrel{DFT}{\leftrightarrow} \quad \underbrace{Y[k]}_{2N-pt} \leftrightarrow \underbrace{C_x[k]}_{N-pt}$$

- this interpretation is useful in various ways
 - it provides insight on why the DCT has better energy compaction
 - it provides a fast algorithm for the computation of the DFT

Energy compaction

$$\underbrace{x[n]}_{N-pt} \leftrightarrow \underbrace{y[n]}_{2N-pt} \quad \stackrel{DFT}{\leftrightarrow} \quad \underbrace{Y[k]}_{2N-pt} \leftrightarrow \underbrace{C_x[k]}_{N-pt}$$

- to understand the energy compaction property
 - we start by considering the sequence y[n] = x[n]+x[2N-1-n]
 - this just consists of adding a mirrored version of x[n] to itself

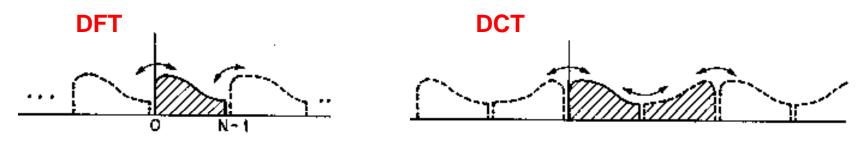


- next we remember that the DFT is identical to the DFS of the periodic extension of the sequence
- let's look at the periodic extensions for the two cases
 - when transform is DFT: we work with extension of x[n]
 - when transform is DCT: we work with extension of y[n]

Energy compaction

$$\underbrace{x[n]}_{N-pt} \leftrightarrow \underbrace{y[n]}_{2N-pt} \quad \stackrel{DFT}{\leftrightarrow} \quad \underbrace{Y[k]}_{2N-pt} \leftrightarrow \underbrace{C_x[k]}_{N-pt}$$

the two extensions are



- note that in the DFT case the extension introduces discontinuities
- this does not happen for the DCT, due to the symmetry of y[n]
- the elimination of this artificial discontinuity, which contains a lot of high frequencies,
- is the reason why the DCT is much more efficient

Fast algorithms

the interpretation of the DCT as

$$\underbrace{x[n]}_{N-pt} \leftrightarrow \underbrace{y[n]}_{2N-pt} \quad \stackrel{DFT}{\leftrightarrow} \quad \underbrace{Y[k]}_{2N-pt} \leftrightarrow \underbrace{C_x[k]}_{N-pt}$$

- also gives us a fast algorithm for its computation
- it consists exactly of the three steps
- 1) y[n] = x[n]+x[2N-1-n]
- 2) Y[k] = DFT{y[n]}this can be computed with a 2N-pt FFT

$$C_{x}[k] = \begin{cases} e^{-j\frac{\pi}{2N}k}Y[k], & 0 \le k < N\\ 0, & otherwise \end{cases}$$

- the complexity of the N-pt DCT is that of the 2N-pt DFT

2D DCT

- the extension to 2D is trivial
- the procedure is the same

$$\underbrace{x[n_1,n_2]}_{N_1 \times N_2 - pt} \leftrightarrow \underbrace{y[n_1,n_2]}_{2N_1 \times 2N_2 - pt} \stackrel{\text{2D DFT}}{\leftrightarrow} \underbrace{Y[k_1,k_2]}_{2N_1 \times 2N_2 - pt} \leftrightarrow \underbrace{C_x[k_1,k_2]}_{N_1 \times N_2 - pt}$$

• with

$$y[n_1, n_2] = x[n_1, n_2] + x[2N_1 - 1 - n_1, n_2] + x[n_1, 2N_2 - 1 - n_2] + x[2N_1 - 1 - n_1, 2N_2 - 1 - n_2]$$

• and

$$C_{x}[k_{1},k_{2}] = \begin{cases} e^{-j\frac{\pi}{2N_{1}}k_{1}}e^{-j\frac{\pi}{2N_{2}}k_{2}}Y[k_{1},k_{2}], & 0 \le k_{1} < N_{1} \\ 0 \le k_{2} < N_{2} \\ 0, & otherwise \end{cases}$$

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2D DCT

• the end result is the 2D DCT pair

$$C_{x}[k_{1},k_{2}] = \begin{cases} \sum_{n_{1}=0}^{N_{1}-1N_{2}-1} 4x[n_{1},n_{2}]\cos\left(\frac{\pi}{2N_{1}}k_{1}(2n_{1}+1)\right)\cos\left(\frac{\pi}{2N_{2}}k_{2}(2n_{2}+1)\right), & 0 \le k_{1} < N_{1} \\ 0 \le k_{2} < N_{2} \\ otherwise \end{cases}$$
$$x[n_{1},n_{2}] = \begin{cases} \frac{1}{N_{1}N_{2}}\sum_{k_{1}=0}^{N_{1}-1N_{2}-1} w_{1}[k_{1}]w_{2}[k_{2}]C_{x}[k_{1},k_{2}]\cos\left(\frac{\pi}{2N_{1}}k_{1}(2n_{1}+1)\right)\cos\left(\frac{\pi}{2N_{2}}k_{2}(2n_{2}+1)\right), & 0 \le n_{1} < N_{1} \\ 0 \le n_{2} < N_{2} \\ 0 & otherwise \end{cases}$$

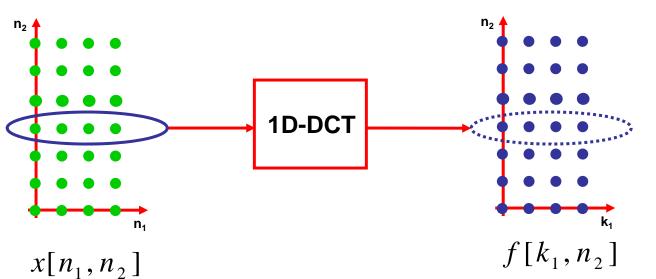
with

$$w_1[k_1] = \begin{cases} \frac{1}{2}, & k_1 = 0\\ 1, & 1 \le k_1 < N_1 \end{cases}, \qquad w_2[k_2] = \begin{cases} \frac{1}{2}, & k_2 = 0\\ 1, & 1 \le k_2 < N_2 \end{cases}$$

• it is possible to show that the 2DCT can be computed with the row-column decomposition (homework)

2D-DCT

 1) create intermediate sequence by computing 1D-DCT of rows



2) compute
1D-DCT of
columns

