

The Fourier Transform & the DFT

- Fourier transform $F(j\omega) = \int_{-\infty}^{+\infty} f(t).e^{-j\omega t}.dt$
- Take N samples of $f(t)$ from 0 to $.(N-1)T$
- Can $F(j\omega)$ be estimated from these?
- Estimate based on rectangular approximation of integral

$$\hat{F}(j\omega) = \sum_{n=0}^{N-1} f(nT).e^{-j\omega nT}.T$$

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- Estimate based on rectangular approximation of integral

$$\hat{F}(j\omega) = \sum_{n=0}^{N-1} f(nT).e^{-j\omega nT}.T$$

- Now take N samples of $\hat{F}(j\omega)$ at multiples of $\omega_0 = \frac{2\pi}{NT}$
- Note that $\hat{F}(j\omega)$ is of period $\frac{2\pi}{T}$

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- Then

$$\hat{F}(jk\omega_0) = \sum_{n=0}^{N-1} T \cdot f(nT) \cdot e^{-j\frac{2\pi}{NT} \cdot k \cdot n \cdot T}$$

$$= T \cdot \sum_{n=0}^{N-1} f(nT) \cdot W_N^{-nk} \quad W_N = \exp\left(j\frac{2\pi}{N}\right)$$

- i.e. DFT relationship where

$$d(n) \stackrel{\Delta}{=} T \cdot f(T \cdot [n \bmod N])$$

- has a DFT $D(k)$ such that

$$D(k) = \hat{F}(jk\omega_0)$$

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- How good is the estimate?

- Let period of DFT operation be $T_0 = NT$

- Hence $T_0 = \frac{2\pi}{\omega_0}$

- Similarly $N\omega_0 = \Omega_0 = \frac{2\pi}{T}$

- There are two approximations involved in estimating $F(j\omega)$

- (i) $f(t)$ is sampled leading to aliasing for non-bandlimited signals.

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- (ii) N samples only are retained

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- Point (i) : Sampling yields a new signal

$$g(n) = T \cdot f(nT)$$

- With Fourier Transform

$$G(j\omega) = \sum_{k=-\infty}^{+\infty} F(j\omega + jk \cdot \Omega_0)$$

- i.e. it may be a poor approximation to $F(j\omega)$
- Point (ii) : Retaining samples $0, N-1$ is effectively windowing the data by

$$w(n) = 1 \quad 0 \leq n \leq N-1$$

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- Fourier Transform of window

$$W(j\omega) = e^{-j \frac{N-1}{2} \cdot \omega T} \cdot \frac{\sin(N\omega T/2)}{\sin(\omega T/2)}$$

- Actual signal used in DFT is $h(n) = g(n) \cdot w(n)$
- i.e. this leads to convolution in frequency domain

$$H(j\omega) = \frac{T}{2\pi} \int_0^{\Omega_0} G(j\phi) \cdot W(j\omega - j\phi) \cdot d\phi$$

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- Note that

$$\begin{aligned}
 H(j\omega) &= \sum_{n=-\infty}^{+\infty} T \cdot f(nT) \cdot w(n) \cdot e^{-j\omega nT} \\
 &= \sum_{n=0}^{N-1} T \cdot f(nT) \cdot w(n) \cdot e^{-j\omega nT} = \hat{F}(j\omega)
 \end{aligned}$$

- I.e. $H(j\omega)$ is the estimate of

$$\hat{F}(j\omega) = \frac{T}{2\pi} \int_0^{\Omega_0} \left[\sum_{k=-\infty}^{+\infty} F(j\phi + jk\Omega_0) \right] \cdot W[j\omega - j\phi] \cdot d\phi$$

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- Since $W(j\omega)$ is periodic with period Ω_0 its contributions to above may be taken over the entire domain as

$$\hat{F}(j\omega) = \frac{1}{2\pi} \int_{-\infty}^{+\infty} F(j\phi) \cdot [T \cdot W[j(\omega - \phi)]] d\phi$$

- i.e. the estimate $\hat{F}(j\omega)$ is the convolution between the desired transform and $V(j\omega)$

$$V(j\omega) = T \cdot e^{-j \frac{(N-1)}{2} \omega T} \cdot \frac{\sin(N\omega T/2)}{\sin(\omega T/2)}$$

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- The main lobe of $V(j\omega)$ is of width $2 \times \frac{2\pi}{NT} = 2\omega_0$
- hence convolution "smears" or "blurs" a step to the same width.
- For greater bandwidth, Ω_0 must be increased (or T decreased).
- To increase resolution we must decrease ω_0
- for a fixed T this implies N must increase.

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- Frequency resolution is the Minimum separation between two sinusoids, resolvable in frequency
- Because of the convolution two impulses in frequency will be smeared, so that if they are to be resolvable they must be separated by at least one frequency bin ie $\frac{2\pi}{NT}$

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- If maximum frequency in the signal is ω_{\max} and the required resolution is Δ then

$$\frac{2\pi}{2T} > \omega_{\max} \quad \Delta > 2 \cdot \frac{2\pi}{NT}$$

- thus $N > \frac{4\pi}{4T} = \frac{4\omega_{\max}}{\Delta}$
- For example (rad/s)

$$\Delta = 2\pi \times 1 \quad \omega_{\max} = 2\pi \times 10^3$$

- then $N > 4000$

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- In practice $\hat{F}(j\omega)$ is required within a prescribed resolution and bandwidth.

- Let $F(j\omega)$ be limited to ω_{\max} , and the required resolution to be Δ or better.

- Then $\Omega_0 > 2\omega_{\max}$ $2\omega_0 < \Delta$

- Hence $\frac{\Omega_0}{2} = \frac{\pi}{T} > \omega_{\max}$ i.e. $T < \frac{\pi}{\omega_{\max}}$

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- I.e. $2\omega_0 = 2 \left[\frac{2\pi}{NT} \right] < \Delta \quad N > \frac{4\omega_{\max}}{\Delta}$
- If $f(t)$ is assumed to be of duration t_{\max} then again $\frac{\Omega_0}{2} > \omega_{\max} \quad \rightarrow \quad T < \frac{\pi}{\omega_{\max}}$

- Set $T_0 = \frac{2\pi}{\omega_0} = NT$

- and $\frac{T_0}{2} > t_{\max} \quad \rightarrow \quad \frac{T_0}{2} = \frac{NT}{2} > t_{\max}$

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- i.e. $N > \frac{2t_{\max}}{T}$

- Hence $N > \frac{2t_{\max} \cdot \omega_{\max}}{\pi}$

where $t_{\max} \cdot \omega_{\max}$ time – bandwidth
product of signal

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Discrete Fourier Transforms

- Consider finite duration signal

$$\{x(n)\} \quad n = 0 \dots N - 1$$

- Its z-transform is

$$X(z) = \sum_{n=0}^{N-1} x(n).z^{-n}$$

- Evaluate at points on z-plane as

$$X(k) = X(z)|_{z_k} = \sum_{n=0}^{N-1} x(n)e^{-j\frac{2\pi}{N}kn}$$

- We can evaluate N independent points

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Discrete Fourier Transforms

$$X(k) = \sum_{n=0}^{N-1} x(n)e^{-j\frac{2\pi}{N}kn}$$

- This is known as the Discrete Fourier Transform (DFT) of $\{x(n)\}$
- Periodic in k ie $X(k + pN) = X(k)$
- This is as expected since the spectrum is periodic in frequency

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Discrete Fourier Transforms

- Multiply both sides of the DFT by $e^{j\frac{2\pi}{N}km}$
- And add over the frequency index k

$$\sum_{k=0}^{N-1} X(k)e^{j\frac{2\pi}{N}km} = \sum_{n=0}^{N-1} x(n) \sum_{k=0}^{N-1} e^{j\frac{2\pi}{N}k(m-n)}$$

- From which $x(n) = \frac{1}{N} \sum_{k=0}^{N-1} X(k).e^{j\frac{2\pi}{N}kn}$

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Discrete Fourier Transforms

- This is the inverse DFT
$$x(n) = \frac{1}{N} \sum_{k=0}^{N-1} X(k).e^{j\frac{2\pi}{N}kn} \quad x(n + qN) = x(n)$$
- That is a) the DFT assumes that we deal with periodic signals in the time domain
b) Sampling in one domain produces periodic behaviour in the other domain

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Discrete Fourier Transforms

- Effectively by knowing $X(z)|_{z=z_k} = X(k)$
- $X(z)$ is known everywhere since

$$X(z) = \sum_{n=0}^{N-1} x(n).z^{-n} = \sum_{n=0}^{N-1} \frac{1}{N} \sum_{k=0}^{N-1} X(k)e^{j\frac{2\pi}{N}kn} z^{-n}$$

- or

$$X(z) = \frac{1}{N} \sum_{k=0}^{N-1} X(k). \frac{1 - z^{-N}}{1 - e^{j\frac{2\pi}{N}k} .z^{-1}}$$

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Discrete Fourier Transforms

- The formula

$$X(z) = \frac{1}{N} \sum_{k=0}^{N-1} X(k). \frac{1 - z^{-N}}{1 - e^{j\frac{2\pi}{N}k} .z^{-1}}$$

- This is essentially an interpolation and forms the basis of the *Frequency Sampling Method* for FIR digital filter design

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Convolution in DFT

- Consider the following transform pairs

$$\{x_p(n)\} \leftrightarrow X_p(k) \quad \{h_p(n)\} \leftrightarrow H_p(k)$$

- Define $Y_p(k) = H_p(k) \cdot X_p(k)$

- Find

$$y_p(n) \stackrel{\Delta}{=} \text{IDFT}[Y_p(k)]$$

Convolution in DFT

- From IDFT

$$y_p(n) = \frac{1}{N} \sum_{k=0}^{N-1} Y_p(k) \cdot e^{j\frac{2\pi}{N}kn} = \frac{1}{N} \sum_{k=0}^{N-1} H_p(k) \cdot X_p(k) \cdot e^{j\frac{2\pi}{N}kn}$$

- However

$$H_p(k) = \sum_{m=0}^{N-1} h_p(m) e^{-j\frac{2\pi}{N}km}$$

$$y_p(n) = \frac{1}{N} \sum_{k=0}^{N-1} \sum_{m=0}^{N-1} h_p(m) \cdot e^{-j\frac{2\pi}{N}km} \cdot X_p(k) \cdot e^{j\frac{2\pi}{N}kn}$$

Convolution in DFT

- Or

$$y_p(n) = \sum_{m=0}^{N-1} h_p(m) \cdot \underbrace{\frac{1}{N} \sum_{k=0}^{N-1} X_p(k) \cdot e^{j\frac{2\pi}{N}k(n-m)}}_{x_p(n-m)}$$

- Thus

$$y_p(n) = \sum_{m=0}^{N-1} h_p(m) \cdot x_p(n-m)$$

- This is the *Circular Convolution*

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Computation of the DFT: The FFT Algorithm

- Computation of DFT requires for every sample N multiplications. There are N samples to be computed i.e. N^2 time consuming operations.
- The Fast Fourier Algorithm: (Decimation in time - DIT, assume even no. of samples)
- set $x_1(n) = x(2n)$ $x_2(n) = x(2n+1)$

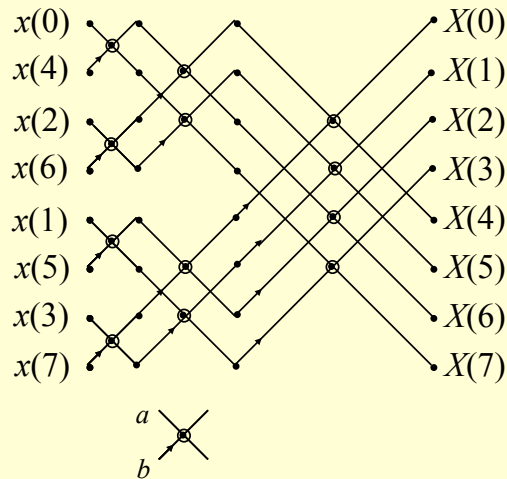
$$n = 0, 1, \dots, \frac{N}{2} - 1$$

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8-point FFT

- 8-point Signal Flow Diagram



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FFT times

- Time (1 multiplication per microsec)

	N	Direct DFT	FFT
2^6	64	.02 sec	.002 sec
2^9	512	1	.02 sec
2^{12}	4096	67	.2
2^{15}	32768	1 hr 11 mins	2
2^{18}	262144	3 days 4 hrs	19

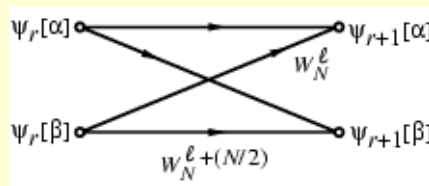
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Decimation-in-Time FFT Algorithm

- In the basic module two output variables are generated by a weighted combination of two input variables as indicated below

where $r = 1, 2, \dots, \mu$ and $\alpha, \beta = 0, 1, \dots, N - 1$



- Basic computational module is called a **butterfly computation**

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Decimation-in-Time FFT Algorithm

- Input-output relations of the basic module are:

$$\Psi_{r+1}[\alpha] = \Psi_r[\alpha] + W_N^\ell \Psi_r[\beta]$$

$$\Psi_{r+1}[\beta] = \Psi_r[\alpha] + W_N^{\ell+(N/2)} \Psi_r[\beta]$$

- Substituting $W_N^{\ell+(N/2)} = -W_N^\ell$ in the second equation given above we get

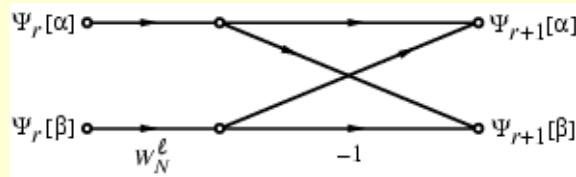
$$\Psi_{r+1}[\beta] = \Psi_r[\alpha] - W_N^\ell \Psi_r[\beta]$$

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Decimation-in-Time FFT Algorithm

- Modified butterfly computation requires only one complex multiplication as indicated below



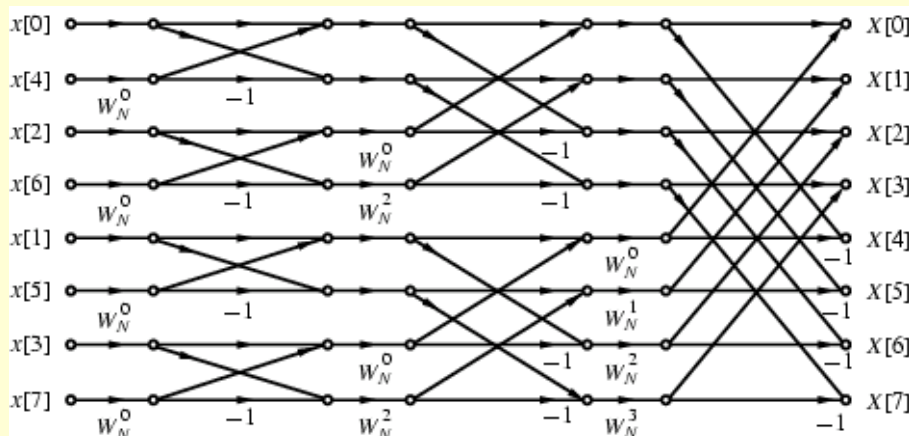
- Use of the above modified butterfly computation module reduces the total number of complex multiplications by 50%

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Decimation-in-Time FFT Algorithm

- New flow-graph using the modified butterfly computational module for $N = 8$



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Decimation-in-Time FFT Algorithm

- Computational complexity can be reduced further by avoiding multiplications by $W_N^0 = 1$, $W_N^{N/2} = -1$, $W_N^{N/4} = j$, and $W_N^{3N/4} = -j$
- The DFT computation algorithm described here also is efficient with regard to memory requirements
- Note: Each stage employs the same butterfly computation to compute $\Psi_{r+1}[\alpha]$ and $\Psi_{r+1}[\beta]$ from $\Psi_r[\alpha]$ and $\Psi_r[\beta]$